SHERLOC: Secure and Holistic Control-Flow Violation Detection on Embedded Systems (Accepted at ACM CCS 2023)

Xi Tan and Ziming Zhao
CactiLab, Department of Computer Science and Engineering, University at Buffalo
{xitan, zimingzh}@buffalo.edu

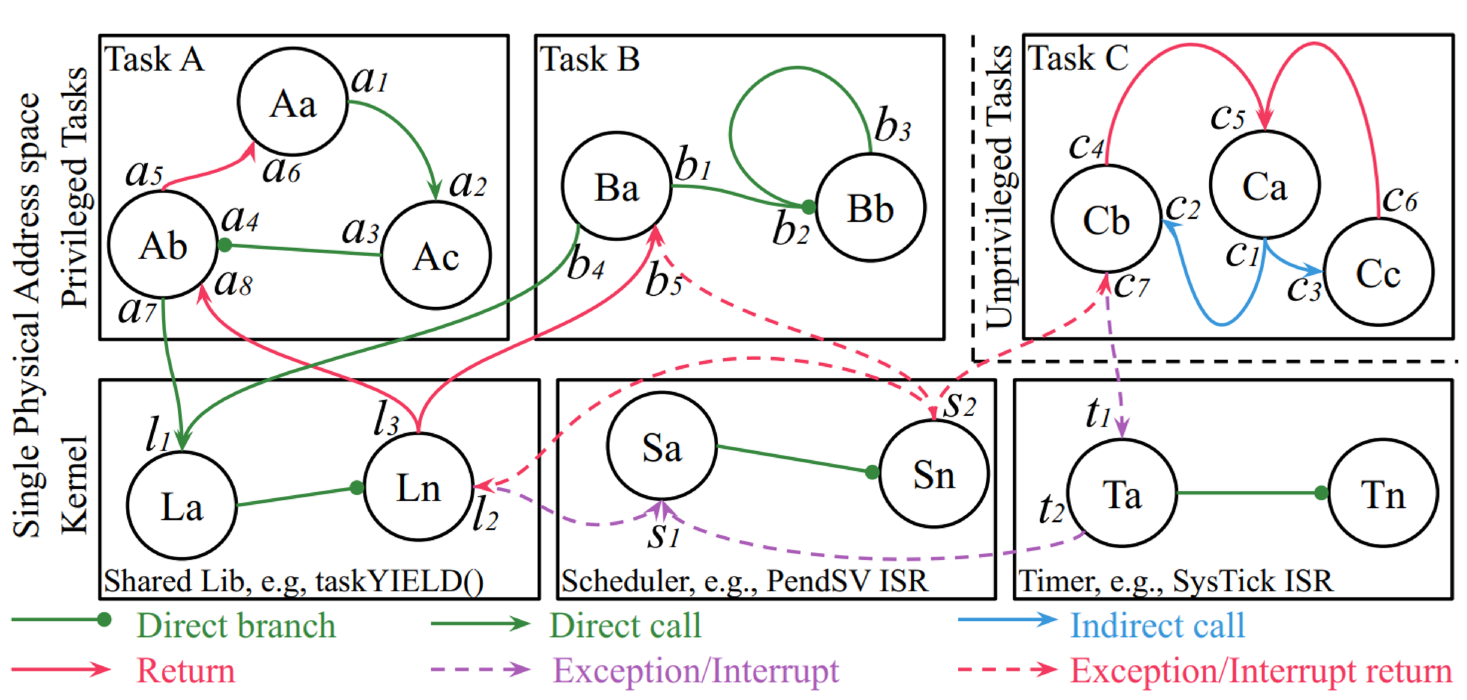
1 Introduction

Microcontroller-based embedded systems are often programmed in low-level languages and are vulnerable to control-flow hijacking. But inlined control-flow integrity (CFI) enforcement solutions increase the binary size and change the memory layout. Trace-based control-flow violation detection (CFVD) offers an alternative, but existing solutions are **application-oriented**, requiring kernel modifications to store and analyze the trace, limiting their use to monitor privileged codes.

2 System-oriented CFVD

Monitor control-flow transfers both within and among privileged and unprivileged components.

- **Interrupt-aware**
  - Interrupts and exceptions occur asynchronously and cannot be anticipated through static analysis or dynamic training. E.g., <c_7, t_1>
- **Scheduling-aware**
  - Scheduler may resume any running tasks' execution. E.g., both $\langle s_2, b_5\rangle$ and $\langle s_2, c_7\rangle$ are legitimate control-flow transfers.
- **Secure hardware tracing**
  - Preventprivileged but potentially compromised syst em to disable tracing
- **Secure trace storage and analysis**
  - Secure trace from the protected system



SCFVD verifies each indirect control-flow transfer must match an edge in the interprocedural CFG (G_S) or the destination must match an address in the interrupt ISR address list (I_K) or the set of task entry or re-entry list (Y_T)

System-oriented CFVD (SCFVD). Given the trace $R_{\mathcal{S}} = (r_0, r_1, \ldots, r_n)$ of a system $\mathcal{S}$ including a kernel $\mathcal{K}$ and tasks $\mathcal{T}$, SCFVD verifies that $r_i \in E_{\mathcal{S}} \vee r_i.d \in I_{\mathcal{K}} \cup Y_{\mathcal{T}}, \forall i \in \{0, 1, \ldots, n\}$.

3 System and Threat Model

- The system features a hardware trace unit. Filtering capabilities are not required.
- The system supports TEE and secure boot for code integrity.
- Attackers can exploit privileged code bugs of the protected system via memory corruptions.

4 Design Overview

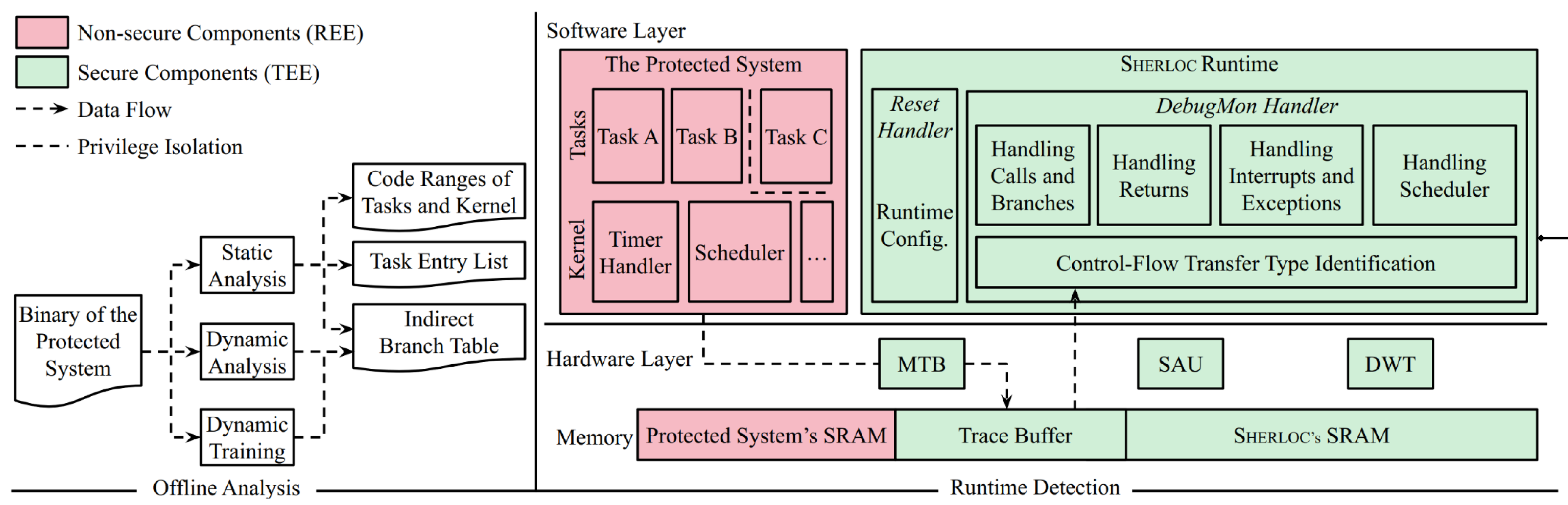


Figure 2: SHERLOC comprises offline analysis and runtime configuration and enforcement modules. The unmodified protected system program runs in the non-secure state, whereas SHERLOC runtime modules execute in the secure state.

5 Runtime Detection Policy

The approach that SHERLOC takes for handling each type of dereferenced instruction in the trace. $\langle s, d\rangle$: a standard trace record. $(\langle s_1, d_1\rangle, \langle s_2, d_2\rangle)$: a pair of interrupt or exception return trace records. IBT: Indirect branch table. VT: non-secure state vector table. RCS: the current task- or kernel-specific reconstructed call stack. Y_T: task entry and re-entry address list.

	Type	Instruction(s)	Ins. Size	How to Identify the Type?	SHERLOC Actions
Bare-metal System and RTOS Cases	Direct branch (§4.4.1)	`B{cond} #imm`	2/4	The dereferenced instruction	Skip the record
	Direct call (§4.4.1)	`BL{cond} #imm`	4	The dereferenced instruction	RCS.push(s + 4)
	Indirect branch (§4.4.1)	`BX{cond} Rx` `TBB/TBH {PC, ...}`	2 4	The dereferenced instruction	*if* $\langle s, d\rangle \notin$ IBT, reset
	Indirect call (§4.4.1)	`BLX Rx`	2	The dereferenced instruction	*if* $\langle s, d\rangle \notin$ IBT, reset; *else* RCS.push(s + 2)

	Function return (§4.4.2)	`BX LR`
		`POP {..., PC}`
		`LDM SP!, {..., PC}`
			2/4	The dereferenced instruction	*if* $d \neq$ RCS.pop(), reset

	Sync. exception (§4.4.3)	`SVC #imm`	2	s[A-bit]	*if* $d \notin$ VT, reset; *else if* $d \neq$ PendSV, RCS.push(s)
	Non-PendSV async. interrupt (§4.4.3)	N/A	N/A	s[A-bit]	*if* $d \notin$ VT, reset; *else if* $d \neq$ PendSV, RCS.push(s)

	Non-PendSV ISR return (§4.4.4)	`BX LR`
		`POP {..., PC}`
		`LDM SP!, {..., PC}`
			2/4	The dereferenced instruction and (d_1 == EXC_RETURN $\wedge$ s_2 == EXC_RETURN)	*if* bare-metal and $d_2 \neq$ RCS.top(), reset; *else if* bare-metal and d_2 == RCS.top(), RCS.pop(); *else* go to PendSV ISR return handling

RTOS-only Cases	PendSV async. interrupt (§4.4.5)	N/A	N/A	s[A-bit]	*if* d == PendSV, $Y_{\mathcal{T}}$.add(s) and $Y_{\mathcal{T}}$.add(RCS.pop())

	PendSV ISR return (§4.4.6)	`BX LR`
		`POP {..., PC}`
		`LDM SP!, {..., PC}`
			2/4	The dereferenced instruction and (d_1 == EXC_RETURN $\wedge$ s_2 == EXC_RETURN)	*if* $d_2 \notin Y_{\mathcal{T}}$, reset; *if* d_2 is in a shared library, and assuming the next trace record is $\langle s_n, d_n\rangle$, and $d_n \notin Y_{\mathcal{T}}$, reset

6 Running Example on FreeRTOS

[] represents RCS with the top on the right-hand side. Black [] represents the active RCS, and gray [] represents an inactive RCS.

Trace Buffer	Runtime Enforcement	RCS for Task A	RCS for Task B	Y_T
$\langle a_3, a_4\rangle$	Direct branch: skip	$[a_1+4]$	[...]	$\{b_4+4, l_2, \ldots\}$
$\langle a_1, a_2\rangle$	Direct call: RCS.push($\boldsymbol{a_1}$ + 4)	$[\boldsymbol{a_1}+4]$	[...]	$\{b_4+4, l_2, \ldots\}$
$\langle a_5, a_6\rangle$	Function return: a_6 == RCS.pop()	[], ~~a_1+4~~	[...]	$\{b_4+4, l_2, \ldots\}$
$\langle l_3, a_8\rangle$	Function return: a_8 is in Y_T ($\boldsymbol{a_8} == \boldsymbol{a_7} + 4$)	$[a_1+4]$	[...]	$\{b_4+4, l_2, \ldots\}$, ~~a_7+4~~
$\langle$EXC_RETURN, $l_2\rangle$ $\langle s_2$, EXC_RETURN$\rangle$	PendSV ISR return: l_2 is in Y_T	$[a_1+4]$	[...]	$\{b_4+4, l_2, a_7+4, \ldots\}$, ~~$l_2$~~
$\langle l_2, s_1\rangle$	s_1 is PendSV ISR address: Y_T.add($\boldsymbol{l_2}$) and Y_T.add($\boldsymbol{b_4}$ + 4)	$[a_1+4]$	[...], ~~b_4+4~~	$\{\boldsymbol{l_2}, \boldsymbol{b_4+4}, l_2, a_7+4, \ldots\}$
$\langle b_4, l_1\rangle$	Direct call: RCS.push($\boldsymbol{b_4}$ + 4)	$[a_1+4]$	$[\ldots, \boldsymbol{b_4}+4]$	$\{l_2, a_7+4, \ldots\}$
$\langle$EXC_RETURN, $b_5\rangle$ $\langle s_2$, EXC_RETURN$\rangle$	PendSV ISR return: b_5 is in Y_T	$[a_1+4]$	[...]	$\{l_2, a_7+4, \ldots\}$, ~~b_5~~
...	...	$[a_1+4]$	[...]	$\{b_5, l_2, a_7+4, \ldots\}$
$\langle l_2, s_1\rangle$	s_1 is PendSV ISR address: Y_T.add($\boldsymbol{l_2}$) and Y_T.add($\boldsymbol{a_7}$ + 4)	$[a_1+4]$, ~~a_7+4~~	[...]	$\{b_5, \boldsymbol{l_2}, \boldsymbol{a_7+4}, \ldots\}$
$\langle a_7, l_1\rangle$	Direct call: RCS.push($\boldsymbol{a_7}$ + 4)	$[a_1+4, \boldsymbol{a_7}+4]$	[...]	$\{b_5, \ldots\}$
$\langle a_3, a_4\rangle$	Direct branch: skip	$[a_1+4]$	[...]	$\{b_5, \ldots\}$
$\langle a_1, a_2\rangle$	Direct call: RCS.push($\boldsymbol{a_1}$ + 4)	$[\boldsymbol{a_1}+4]$	[...]	$\{b_5, \ldots\}$
...	...	[]	[...]	$\{b_5, \ldots\}$

Direct transfer | Function return | ISR entry | ISR return

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CactiLab   Code Repo

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